



# Polynomial DC decompositions

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### Difference of convex (dc) programming

#### Problemss of fthe form

$$\min f_0(x)$$
  
s. t.  $f_i(x) \le 0, i = 1, ..., m$ 

where  $f_i(x) \coloneqq g_i(x) - h_i(x)$ ,  $g_i$ ,  $h_i$  convex for i = 0, ..., m. where , convex for

#### What if such a decomposition is not given?

regularization) procedures of some kind. However, all the proofs we know are "constructive" in the sense that they indeed yield g and h satisfying (2.3) but could hardly be carried over computational aspects.

Hiriart-Urruty, 1985

stance. For one thing, the d.c. structure of a given problem is not always apparent or easy to disclose, and even when it is known explicitly, there remains for the problem solver the hard task of bringing this structure to a form amenable to computational analysis. However, since the d.c. structure is virtually universal, one can hope that

Tuy, 1995





### Difference of convex (dc) decomposition

• Difference of convex (tt) decomposition: given a polynomial, finish d aga can soluth sthoant that

$$f=g-h$$
,

where *g, h* convex polynomials. where convex polynomials.

- Questions:
   Questions:
   Does such a decomposition always exist?
  - Pear suchanderemposition always existing?

  - Is this decomposition unique?





### Existence of dc decomposition (1/4)

\*•Applynomial is a summof square (sqs) if, polynomial is a summof square (sqs) if, polynomial is a summof square (sqs) if  $p(x) = \sum q_i^2(x)$ .

• A polynomial p(x) of degree 2d is sos if and only if  $\exists Q \geq 0$  such that  $p(x) = z(x)^T Q z(x)$ 

where  $z=[1,x_1,\dots,x_n,x_1x_2,\dots,x_n^d]^{\mathrm{T}}$  is the vector of monomials up to where  $z=[1,x_1,\dots,x_n,x_1x_2,\dots,x_n^d]^{\mathrm{T}}$ 

•• Testing whether a polynomial is sos is a semidefinite program.





### Existence of dc decomposition (2/4)

$$\begin{array}{c}
f(x) \\
\text{convex}
\end{array} \Leftrightarrow \begin{array}{c}
H_f(x) \geqslant 0, \\
\forall x \in \mathbb{R}^n
\end{array} \Leftrightarrow \begin{array}{c}
y^T H_f(x) y \ge 0, \\
\forall x, y \in \mathbb{R}^n
\end{array} \Leftarrow \begin{array}{c}
y^T H_f(x) y \text{ sos}
\end{array}$$
Sos-convexity

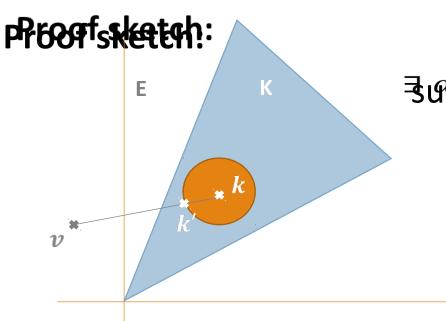
Theorem: Any polynomial can be written as the difference of two sos-convex polynomials.

Corollary: Any polynomial can be written as the difference of two convex polynomials.





# Existence of dc decomposition (3/4)



$$=: k'$$

$$\exists \text{ such that } (1-\alpha)v + \alpha k \in K$$

$$\Leftrightarrow v = \frac{1}{1 - \alpha} k' - \frac{\alpha}{1 - \alpha} k$$

$$k_1 \in K \qquad k_2 \in K$$





# Existence of dc decomposition (4/4)

- ••Heere | PO | \tau | pooligitoon idegroe degriee 201, imples ariables },

  sks-co | sues-pohyreompielly pooleigite eo 2 dlag de ien 2 dvaniabiles variables }.
- ••Remains to how what at Kulisdichled incentional:

  (Ser pe showe to be in the tipe of terior of K.

• Also shows that a decomposition can be obtained efficiently:

solving 
$$f = g - h$$
, is an SDP.

• In fact, we show that a decomposition can be found via LP and SOCP (not covered here).

output

output

not covered here).





### Uniqueness of dc decomposition

•• Decemposition: given a polynomial , f in the general k examples spectral batials such that f = g - h.

#### ·· Questions:

- •• Doesssuch adecomposition always exist? Ves
- •• Can lobtain such a decomposition efficiently? ✓ Through sos-convexity
- Is this decomposition unique?

#### Initial decomposition

$$f(x) = g(x) - h(x)$$



#### Alternative decompositions

$$f(x) = (g(x) + p(x)) - (h(x) + p(x))$$
$$p(x) \text{ convex}$$

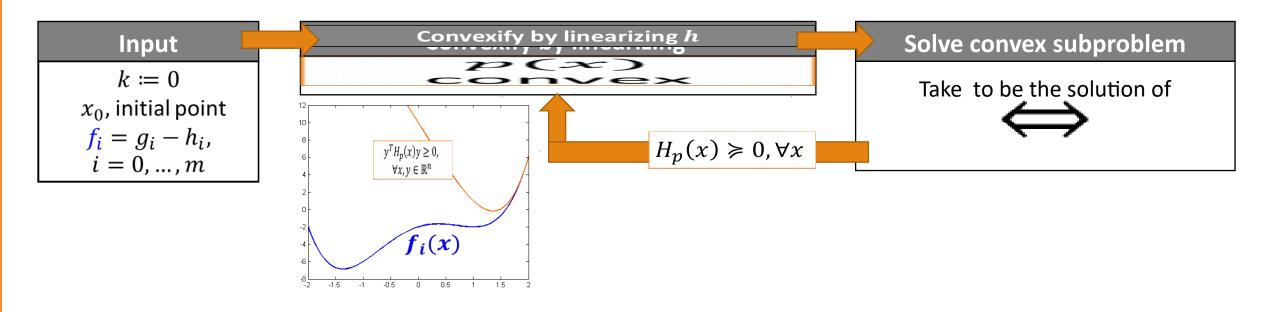
"Best decomposition?"





### Convex-Concave Procedure (CCP)

- Heuristic for minimizing DC programming problems.
- Idea:

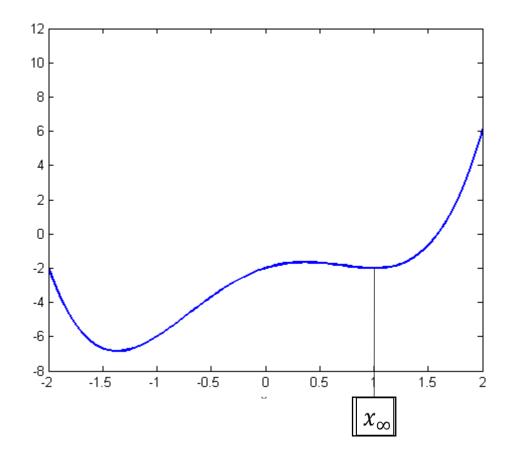






### Convex-Concave Procedure (CCP)

•• Troyyexamples:,  $\underset{x}{\text{wine}} f \in x$ , ....



Reiterate





### Picking the "best" decomposition for CCP

#### Algorithm

Linearize h(x) around a point  $x_k$  to obtain convexified version of f(x)

#### Idea

Pick h(x) such that it is as close as possible to affine around  $x_k$ 

#### **Mathematical translation**

Minimize curvature of h at  $x_k$ 

#### **Worst-case curvature\***

$$\min_{\mathbf{g},\mathbf{h}} \lambda_{max}(H_h(x_k))$$
s.t.  $f = g - h$ 
 $g, h \text{ convex}$ 

$$*\lambda_{max}H_h(x_k) = \max_{y \in S^{n-1}} y^T H_h(x_k) y$$

#### Average curvature\*

$$\min_{g,h} Tr H_h(x_k)$$
s.t.  $f = g - h$ ,
 $g, h$  convex

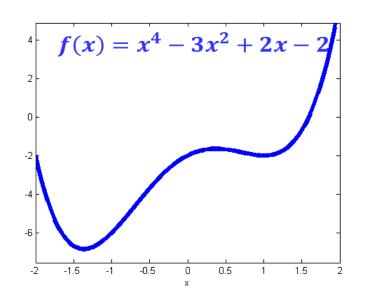
\* 
$$Tr H_h(x_k) = \int_{y \in S_{n-1}} y^T H_h(x_k) y d\sigma$$





# Undominated decompositions (1/4)

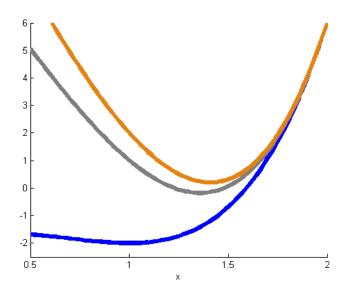
Definition: (g, h := g - f) is an undominated decomposition of f if no other decomposition of f can be obtained by subtracting a (nonaffine) convex function from g.





#### **DOMINATED BY**

$$g'^{(x)}=x^4, \ h'(x)=3x^2+2x-2$$
 Convexify around  $x_0=2$  to Convexify around to get







# Undominated decompositions (2/4)

$$D = \begin{pmatrix} 0 & d_{12} & * & d_{14} \\ d_{12} & 0 & d_{23} & * \\ * & d_{23} & 0 & d_{34} \\ d_{14} & * & d_{34} & 0 \end{pmatrix}$$

$$\mathbf{x_1} = \begin{pmatrix} x_1^1 \\ x_1^2 \\ x_1^3 \end{pmatrix}, \mathbf{x_2} = \begin{pmatrix} x_2^1 \\ x_2^2 \\ x_2^3 \end{pmatrix}, \dots$$

Incomplete distance matrix



Rerever dotation by the points in  $\mathbb{R}^d$ 

Solve:

$$\min_{x_{i}, x_{j} \in \mathbb{R}^{d}} \sum_{i < j} (\|x_{i} - x_{j}\|_{2}^{2} - d_{ij}^{2})^{2}$$

There is a realization in  $\mathbb{R}^d$  of the value = 0.





# Undominated decompositions (3/4)

$$\min_{x_{i},x_{j}} \sum_{i < j} (\|x_{i} - x_{j}\|_{2}^{2} - d_{ij}^{2})^{2}$$

$$= \|x_{i} - x_{j}\|_{2}^{4} + d_{ij}^{4} - 2d_{ij}^{2} \cdot \|x_{i} - x_{j}\|_{2}^{2}$$

$$g_{ij} \qquad h_{ij}$$

(is an undominated acd of the objective function, dcd of the objective function.





# Undominated decompositions (4/4)

Theorem: Giverna podynomial, considerder

Any aptimalisely tion is an underninated ded of partial application as what so exists).

Theorem: If fast deglesser et, ist strong typhy Phrahama to to told ve (\*).

Replace  $f_0 = g_0 + h_0 + g_0 + h_0 + g_0 + h_0 + g_0 + g_$ 





### Comparing different decompositions (1/2)

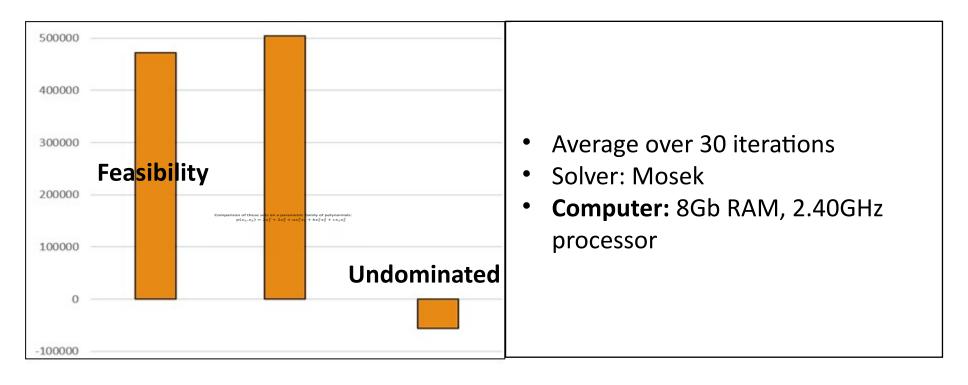
••Solwingtheproblem:, where has findwhere  $f_0$  has n=8 and d=4.
• Decompose run CCP for 4 minutes and compare objective value.
• Decompose  $f_0$ , run CCP for 4 minutes and compare objective value.

Feasibility	Feasibility	$\lambda_{max}H_h(x_0)$	Undominated	Undominated
$\min_{\substack{g,h\\\text{s.t.}}} 0$ s.t. $f_0^{S} = \frac{t}{g} \cdot g - h$ $g$ , $SOSOSOMWEX$	•		$t$ $g - h$ $vex$ $onvex$ $_0) \ge 0$	$\min_{g,h} \frac{1}{A_n} \int_{S_{n-1}} Tr \ H_g \ d\sigma$ $s. \cot_{\sigma} \cot_{\sigma} H_g $





### Comparing different decompositions (2/2)



**Conclusion:** Rate of convergence of CCP strongly affected by initial decomposition.





### Main messages

- We studied the question of decomposing a polynomial into the difference of two convex polynomials.
- This decomposition always exists and is not unique.
- Choice of decomposition can impact convergence rate of the CCP algorithm.
- Dc decompositions can be efficiently obtained using the notion of sos-convexity (SDP)
- Also possible to use LP or SOCP-based relaxations to obtain dc decompositions (not covered here).





# Thank you for listening

Questions?

Want to learn more? <a href="http://scholar.princeton.edu/ghall">http://scholar.princeton.edu/ghall</a>

7/31/16