

# IP MULTICAST

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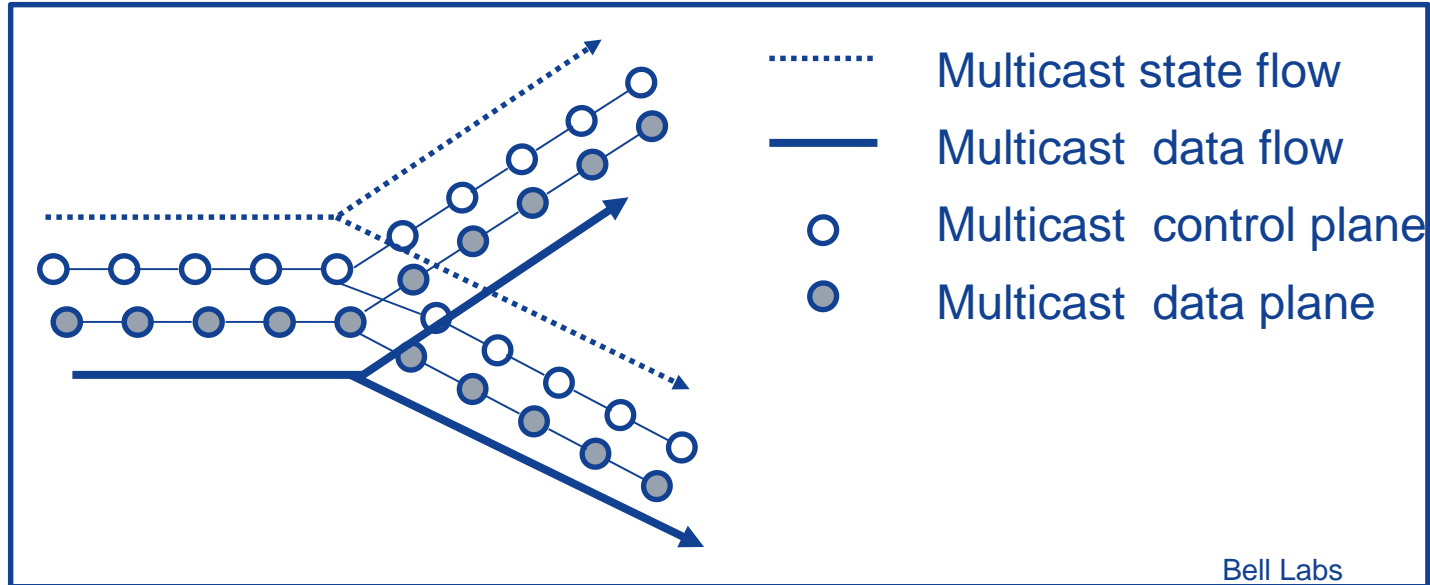
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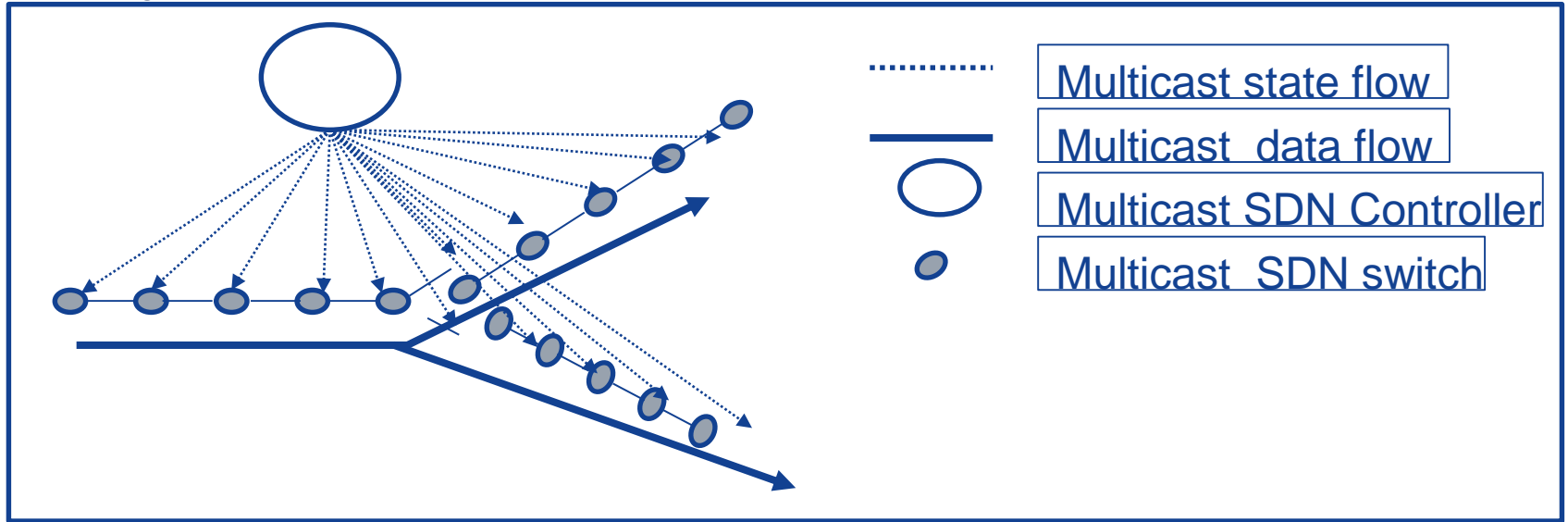
# Why is IP Multicast not deployed in public networks?

- Denial-of-Service (DoS) attack amplification
- Complex Control Plane
- Large Forwarding state
  - Non aggregable



# Can SDN help with Multicast?

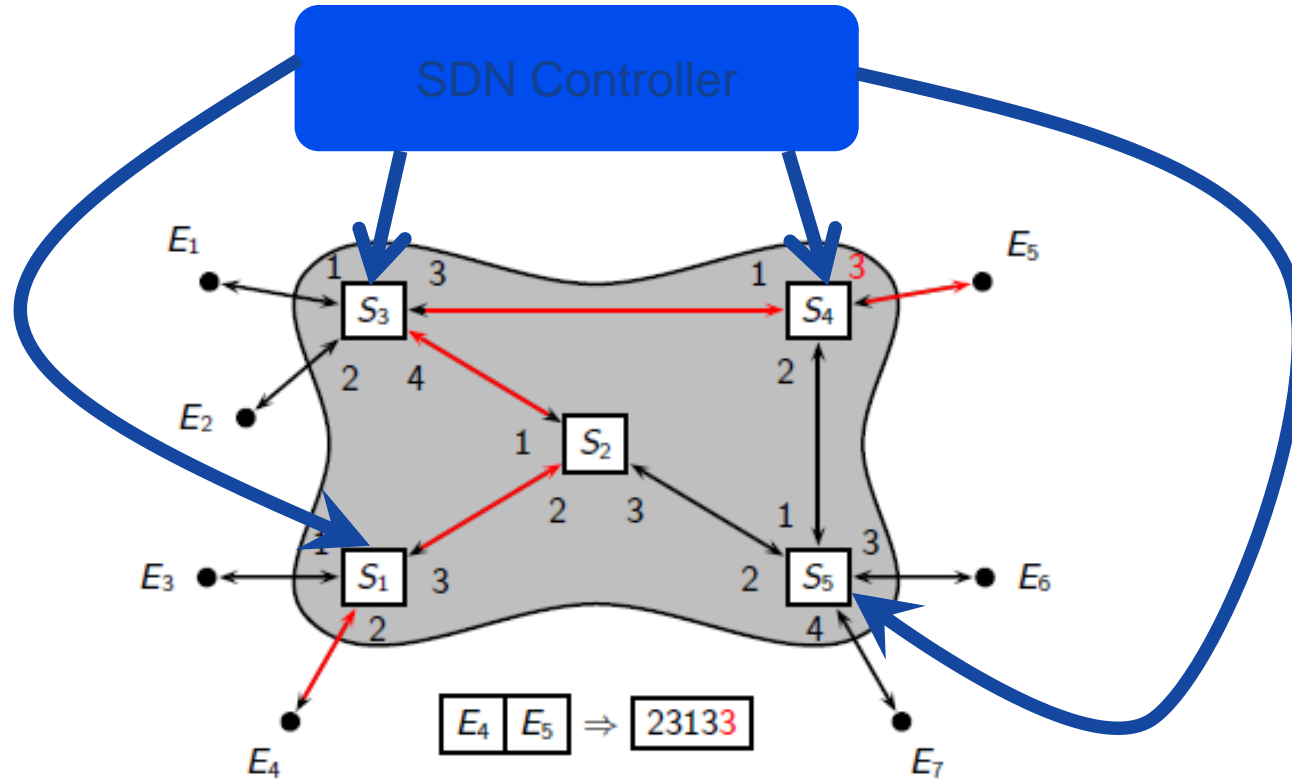
- Denial-of-Service (DoS) attack amplification ✓
- Control state ✓
- Forwarding state ✗



# Can we eliminate multicast forwarding state in SDN?

- Eliminate unicast forwarding state in SDN:
  - Path Switching: **per-flow routing without per-flow state**
  - New data path suitable for SW switches and programmable packet processors
  - Encode path in the packet headers
  - DIMACS 2016

# Eliminating unicast forwarding state in SDN using Path Switching



# Can we eliminate multicast forwarding state in SDN?

Can we extend Path Switching to encode multicast paths?

Can we create an efficient encoding of a multicast path?

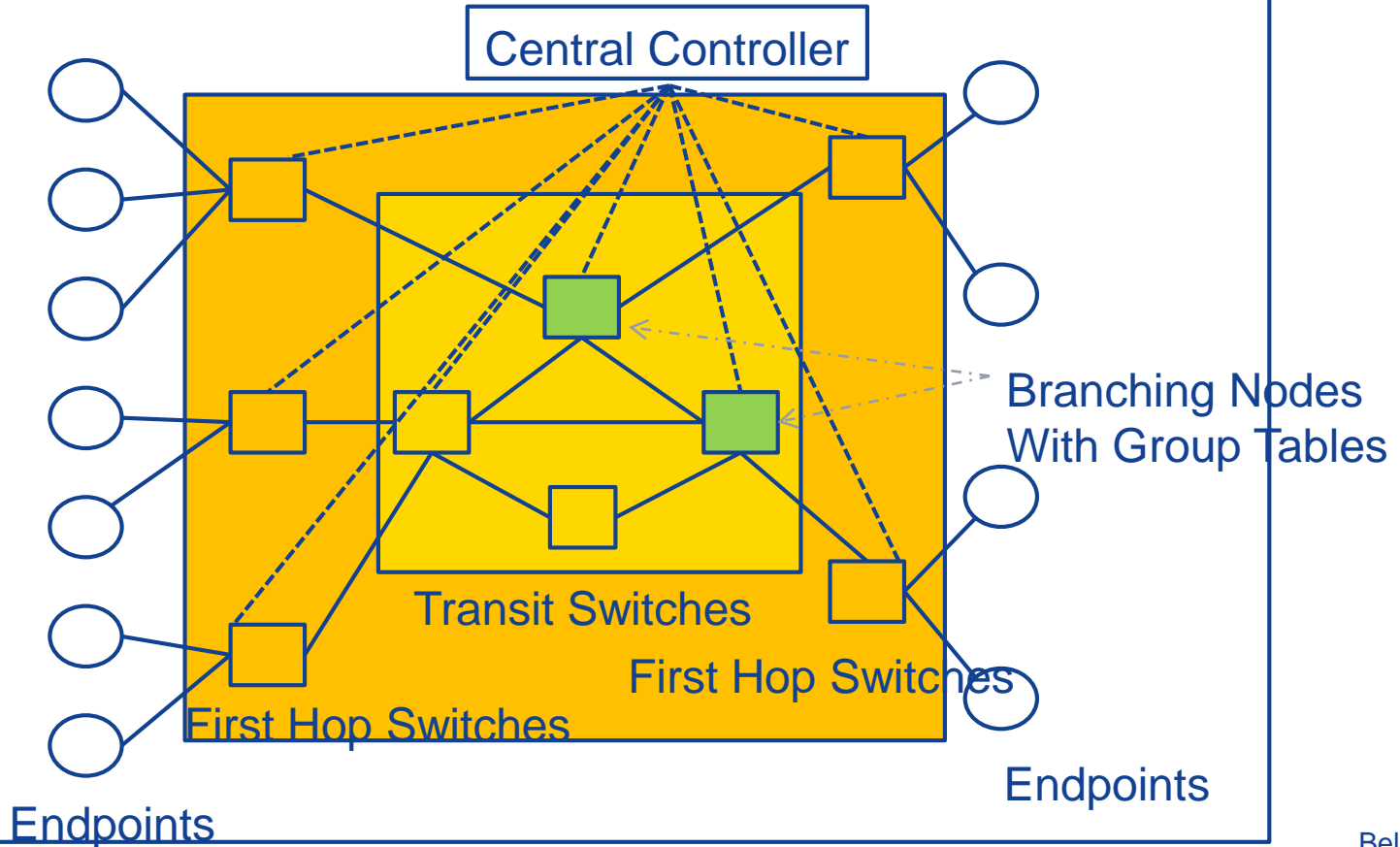
- No blowup in packet size (e.g. using bitmaps)

- No blowup in storage state (e.g., encode each multicast tree by a unique identifier)

# Can we reduce multicast forwarding state in SDN?

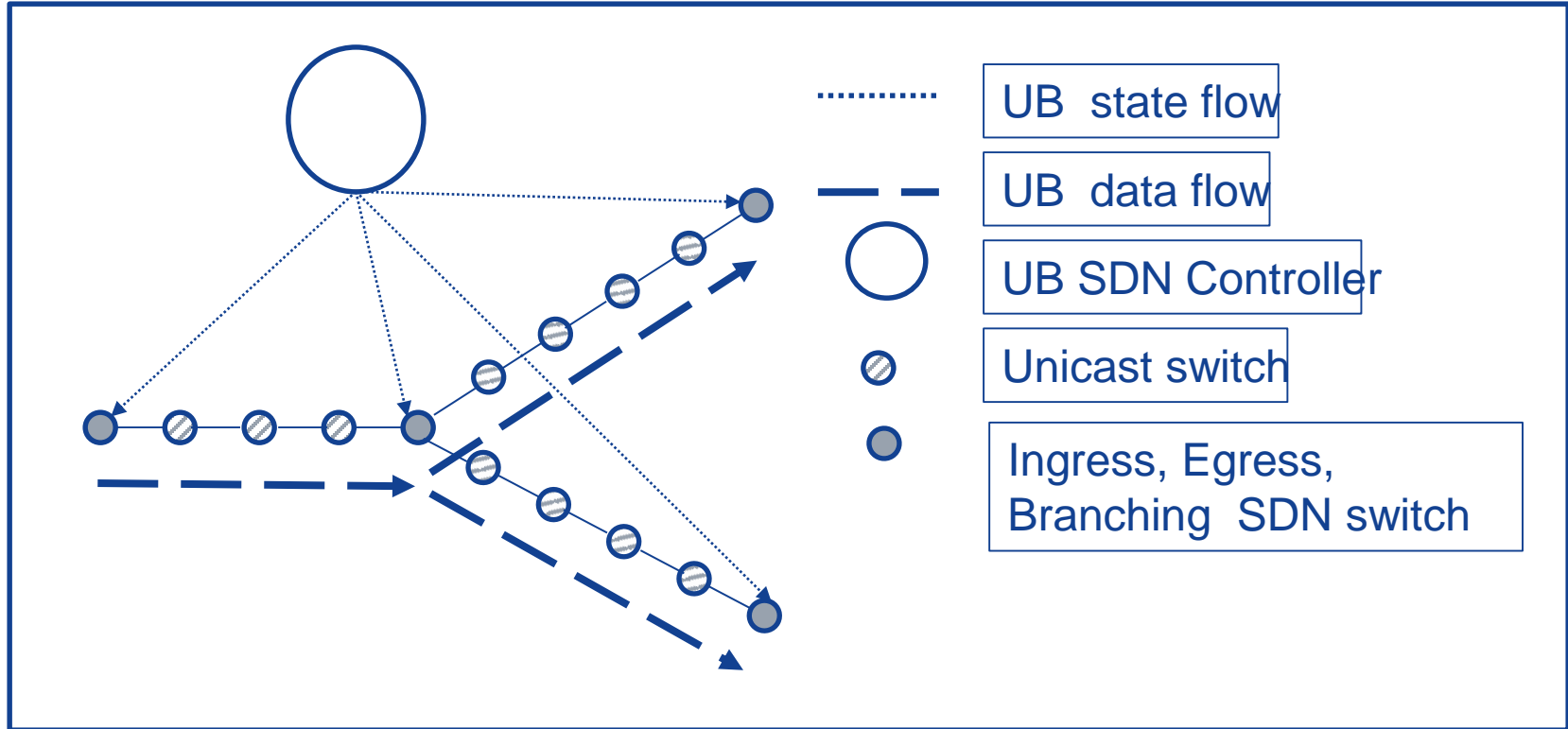
- Unicast Branching (UB)
  - Use branching nodes in the network to replicate unicast flows.
  - Use SDN Flow Table at ingress and egress
  - Use SDN Group Table at branching nodes
- Reduces multicast forwarding state from linear to sublinear in number of forwarding nodes

# Unicast Branching (UB) Reference Diagram



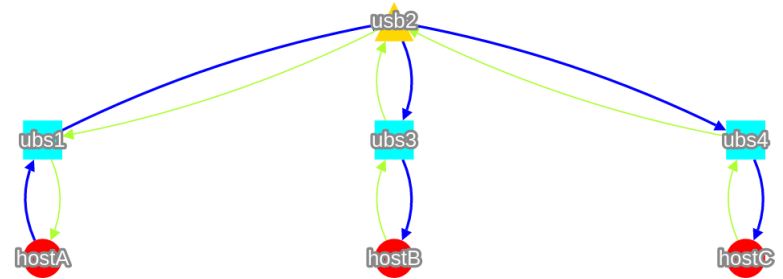


# Reducing multicast forwarding state in SDN using Unicast Branching (UB)



# Added advantages of Unicast Branching (UB)

- **Tunable knob** to switch between unicast replication and full multicast
- Allows for an **NFV** based implementation
- Allows **Traffic Engineered** branches
  - Fast Reroute, Per branch QoS
- Works at all protocol layers – protocol agnostic
  - Ethernet, IP, MPLS
- **Enables unicast only protocols like Segment Routing and TCP to be multicast capable\***
  - HTTP Adaptive Streaming multicast
  - Efficient content caches
- Enables **Policy Based Multicast**



\* Requires stateful NFV elements, not just SDN switches for branching points

# Policy Based Multicast

- Policy based networking: Rules for non default routing
  - Geofencing
  - QoS
  - Membership filtering
- UB enables Policy Based Multicast
  - Number, location and type of branching nodes

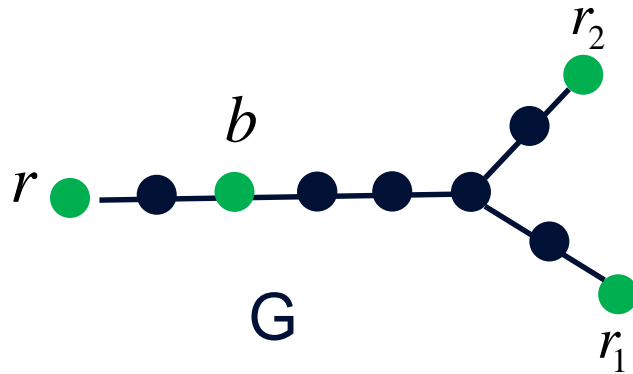
# Where are the Algorithms?

# Building Efficient Policy Based Multicast Trees

- Problem 1 definition:
  - Given an ingress node, a set of egress nodes and a set of branching nodes, build an “optimal” multicast tree.
  - What is “optimal”
  - Usual definition is based on link cost.
  - Steiner tree problem (NP-complete)

# Building Multicast Trees using UB – Major Issue

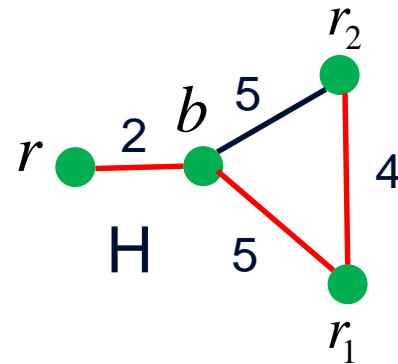
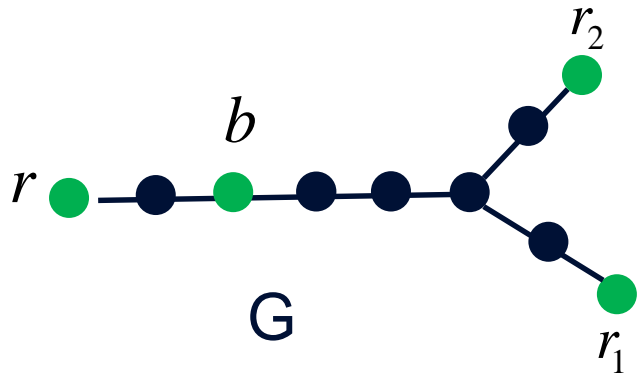
- UB based multicast tree is not a tree!!!
  - It is a “configuration”
- Cannot apply Steiner tree approximation solutions directly.
- Problem: How to create minimum cost configurations?



# Transformation to Steiner tree problem on H

Define :

- Edge-weighted graph  $H = (O, E)$ .  $O$  is set of branching nodes (including terminals)
- $e=(b,b') \in O$ ,  $w(e) = \text{length shortest path containing no internal } O \text{ nodes}$



**Theorem:** Minimum cost configuration problem in  $G$  is equivalent to Steiner tree problem in  $H$

# Minimum cost configuration problem

**Theorem:** There is a polynomial-time 1.39-approximation algorithm for min cost configuration problem. [BGRS10]

**Theorem:** The minimum cost configuration problem is APX-hard.

**Proof:** Follows from APX-hardness of Steiner problem for complete graphs with weights 1 and 2. [BP89]



## Problem 2: Minimize branching nodes

- Problem 1: Minimize cost given a set of branching nodes . MIN COST PROBLEM
- Problem 2: Minimize number of branching nodes given a fixed cost. MIN BRANCHING PROBLEM

## Min Branching Problem

- For a subset  $X$  of the transit nodes, let  $C_x$  be the minimum cost valid configuration using  $X$  as the set of extra branching nodes.
- We are given a graph  $G = (V, E)$ , a multicast demand  $d = (r, r_1, r_2, \dots, r_t)$ , a bound  $k$  and an attainable cost  $c$ .
- Does there exist a branching set  $X$  with least cost valid configuration  $C_x$  satisfying  $d$  where  $|X| \leq k$  and  $\text{cost}(C_x) \leq c$ .

# Min Branching Problem

**Theorem:** This problem is NP-complete.

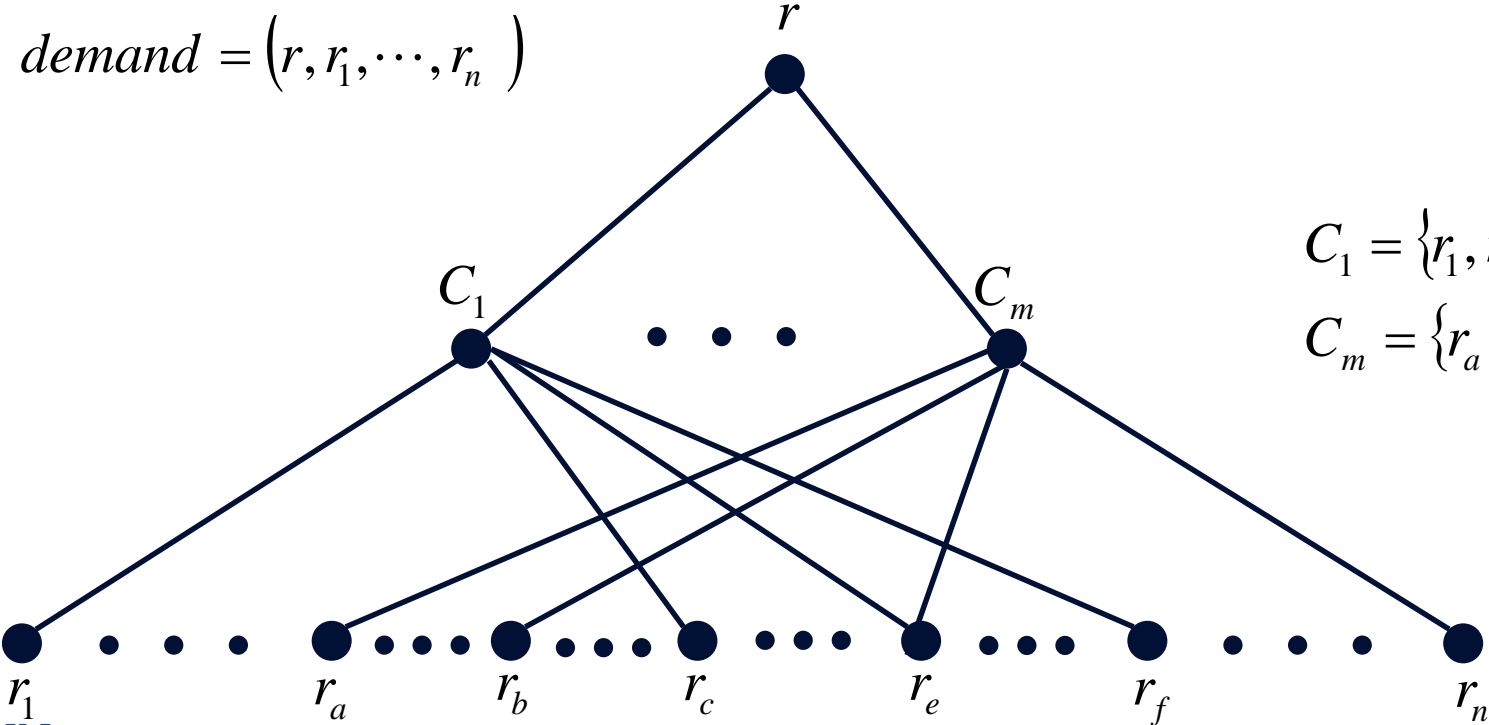
**Proof:** Follows from a construction using Set Cover. ■

**Corollary:** For this problem the best possible approximation is  $\approx \ln n$ .

**Proof:** Follows from bounds for Set Cover. ■

# Theorem: Min Branching is NP-complete

$$\text{demand} = (r, r_1, \dots, r_n)$$



$$C_1 = \{r_1, r_c, r_e, r_f\}$$

$$C_m = \{r_a, r_b, r_e, r_n\}$$

# Policy Driven Software Defined Multicast Using Efficient Selection of Unicast Branching Points

- Conclusion:
  - Unicast Branching based multicast provides for efficient, policy driven Software Defined Multicast.