

THE DISCREPANCY OF  
QUASI-ARITHMETIC PROGRESSIONS:  
POWER OF  $N$  OR POWER OF  $\log N$ ?

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Van der Waerden's Theorem: For all  $k$ , there exists  $N=N(k)$  such that every 2-colouring of  $\{1, 2, \dots, N\}$  yields a monochromatic  $k$ -term arithmetic progression.

Bounds: For all colourings,  $k \geq \log \log \log \log \log N$  (Gowers)  
There exists a colouring with  $k \leq \log_2 N$  (Local Lemma)

Roth's Theorem: For all  $k$ , there exists  $N=N(k)$  such that every 2-colouring of  $\{1, 2, \dots, N\}$  yields an arithmetic progression of discrepancy at least  $k$ .

Discrepancy = |Number of reds - Number of blues|

Bounds: For all colourings,  $k \geq c_1 N^{1/4}$  (Roth)  
There exists a colouring with  $k \leq c_2 N^{1/4}$  (Matoušek & Spencer)

[Sarkozy:  $N^{1/3+\epsilon}$ , Beck:  $N^{1/4+\epsilon}$ ]

Homogeneous Arithmetic Progressions:  $\{0, d, 2d, 3d, \dots\}$

Is there a colouring of  $\mathbb{N}$  such that all homogeneous arithmetic progressions have bounded discrepancy? (Erdős)

Upper Bound:  $O(\log n) \rightarrow \chi(3^a b) = \begin{cases} 1 & \text{if } b \equiv 1 \pmod{3} \\ -1 & \text{if } b \equiv -1 \pmod{3} \end{cases}$

Partial Colouring:  $\chi(3k+1) = 1 \quad \chi(3k) = 0 \quad \chi(3k-1) = -1$

Upper bound in terms of  $d$  alone:  $O(d^{4+\epsilon})$  (Reimer)

Quasi-arithmetic Progressions:  $\{0, [\alpha], [2\alpha], [3\alpha], \dots\}$

e.g:  $\alpha = \sqrt{5} \rightarrow \{0, 2, 4, 6, 8, 11, 13, \dots\}$

There are  $O(N^2)$  distinct quasi-arithmetic progressions contained in  $\{0, 1, 2, \dots, N\}$  (Farey Sequence)

Polynomial-size family, random colouring upper bound on discrepancy:  $O(\sqrt{N \log N})$

## Lower Bounds on the discrepancy of quasi-progressions

For any 2-colouring of  $\{0, 1, \dots, N\}$ , there exists a quasi-progression of discrepancy at least  $c(\log N)^{1/4}$   
(Hochberg '94)

Improvement: For any 2-colouring of  $\{0, 1, \dots, N\}$ , there exists a quasi-progression of discrepancy at least  $\frac{1}{50} N^{1/6}$

Proof Sketch:  $\rightarrow$  Pick a subinterval  $\{N-M, N-M+1, \dots, N\}$   
 $\rightarrow$  Take the A.P. of discrepancy  $\frac{1}{50} M^{1/4}$  and common difference  $d$  (fact:  $d \leq \sqrt{6M}$ ) inside  $\{N-M, N-M+1, \dots, N\}$   
 $\rightarrow$  Show that this A.P. can be realised as a quasi-progression.

$$\sqrt{8} \rightarrow \{0, \underbrace{2, 5, 8, 11, 14}_{\equiv 2 \pmod{3}}, \underbrace{16, 19, 22, 25, 28}_{\equiv 1 \pmod{3}}, \dots\}$$

$$\alpha = d - \epsilon$$

$\lfloor k\alpha \rfloor \equiv -1 \pmod{d}$  for the first  $(1/\epsilon)$  terms  
 $\equiv -2 \pmod{d}$  for the next  $(1/\epsilon)$  terms  
 and so on.

For  $M = N^{2/3}$ , all arithmetic progressions inside  $\{N-M, N-M+1, \dots, N\}$  can be realised as quasi-progressions.

## Typical Behaviour of Quasi-progressions

Given any 2-colouring of the non-negative integers, for almost every  $\alpha \in [1, \infty)$ , there are infinitely many  $n$  such that  $D_\alpha(n) \geq \log^* n$  (Beck '86)

$$D_\alpha(n) = \max_{1 \leq k \leq n} \left| \sum_{j=0}^k \chi(\lfloor \alpha j \rfloor) \right|$$

Proof Sketch: Let  $N_t(k)$  be the smallest integer such that any 2-colouring of an interval  $J$  of length  $N_t(k)$  yields, for "most"  $\alpha \in [t, t+1)$ ,  $D_\alpha(j) \geq k$  for some  $j \in J$

Now consider  $2^{N_t(k)}$  blocks of length  $N_t(k)$ . There must be two blocks coloured identically

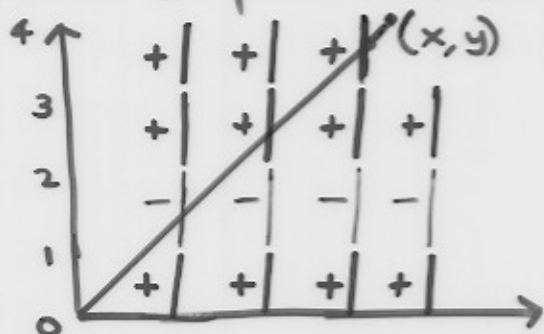
$$\begin{array}{cc} \overline{a_1 \quad b_1} & \overline{a_2 \quad b_2} \\ \mathcal{D}_\alpha(j) \geq k & \mathcal{D}_\alpha(j) \geq k \\ \hline L: \text{Number of terms in } [a_1, a_2] \approx \frac{a_2 - a_1}{\alpha} \end{array}$$

If  $L$  is odd, either  $[b_1, a_2]$  or  $[a_1, b_2]$  yields discrepancy at least  $k+1$ .

Thus  $N_t(k+1) \leq 2^{N_t(k)}$  i.e.,  $D_\alpha(n) \geq \log^* n$

Improvement: Given any 2-colouring of the non-negative integers, for almost every  $\alpha \in [1, \infty)$ , there are infinitely many  $n$  such that  $D_\alpha(n) \geq (\log n)^{3/5}$

(Also works for partial colourings of positive density)

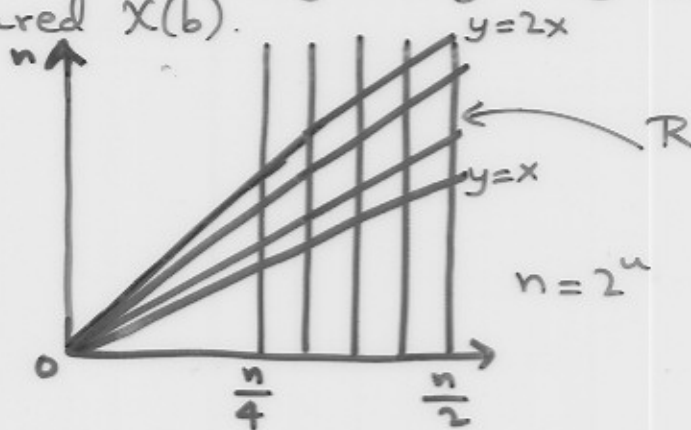


$$\alpha \in [1, 2)$$

$$D(x, y) = \sum_{k \leq x} \chi(\lfloor k\alpha \rfloor)$$

$$\alpha = y/x$$

The vertical line segment joining  $(a, b)$  with  $(a, b+1)$  is coloured  $\chi(b)$ .



Defn:  $\alpha$  is balanced iff  $D_\alpha(n) < (\log n)^{3/5}$

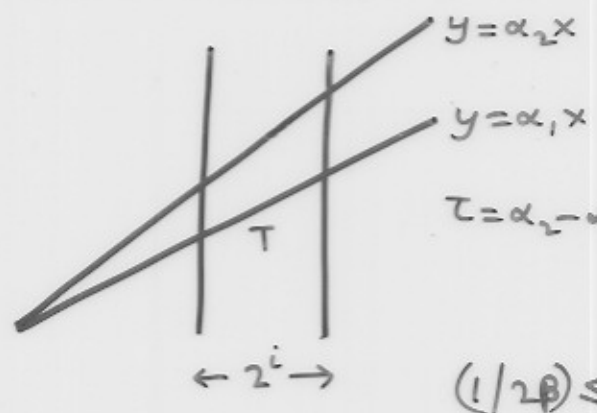
$$H(x, y) = \begin{cases} D(x, y) & \text{if } y/x \text{ is balanced} \\ 0 & \text{otherwise} \end{cases}$$

Plan: Suppose the set of balanced  $\alpha$  has measure  $\delta > 0$ .

Construct orthonormal functions  $g_1, g_2, \dots, g_r$  where  $r = \frac{\log n}{8}$  and  $\sum_{i=1}^r \langle H, g_i \rangle^2 \geq c_0 n^2 (\log n)^{3/5}$

Since  $R$  has area  $O(n^2)$ , Bessel's inequality yields a contradiction.

The  $i^{\text{th}}$  trapezoidal grid



$\tau = \alpha_2 - \alpha_1 = \frac{1}{2^i n^\beta}$  ( $\beta$  to be fixed later; around  $(\log n)^{1/5}$ )

$(1/2\beta) \leq \text{area of } T \leq (1/\beta)$

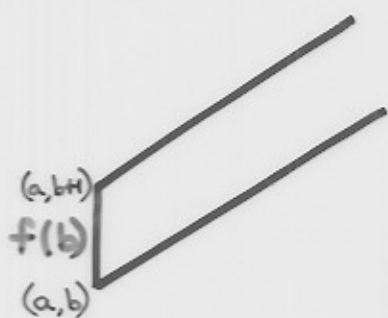
- The position of the leftmost vertical line is chosen randomly and uniformly in  $(\frac{n}{4}, \frac{n}{4} + 2^i)$
- The slope of the lowermost slanting line is chosen randomly and uniformly in  $(1, 1 + \tau)$

Switch values:  $\{b \mid X(b) \neq X(b-1)\}$

Switch points:  $\{(a, b) \mid b \text{ is a switch value}\}$

Observation: If  $X$  is constant on some interval of length  $2(\log n)^{1/5}$ , then for all  $\alpha \in [1, 2)$ , we have  $D_\alpha(n) \geq (\log n)^{1/5}$ .

If not, there are at least  $\frac{n}{2(\log n)^{1/5}}$  switch values and  $\frac{n^2}{4(\log n)^{1/5}}$  switch points.



$$H_{a,b}(x,y) = \begin{cases} f(b), & \text{if } \frac{b}{a} \leq \frac{y}{x} < \frac{b+1}{a} \text{ and } \\ & \frac{xy}{a} \text{ is balanced} \\ 0 & \text{otherwise.} \end{cases}$$

$$H(x,y) = \sum_{a=0}^{\infty} \sum_{b=0}^{\infty} H_{a,b}(x,y) \quad [\text{Finite sum for fixed } (x,y)]$$

Good switch point: Does not share its trapezoid with another lattice point, for any positioning of the grid.

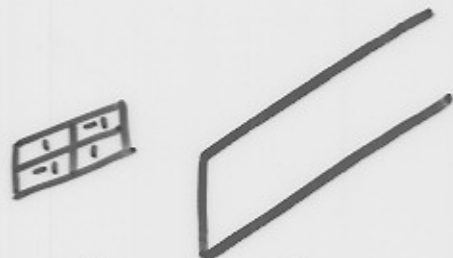
$$G_i = \begin{cases} \begin{array}{|c|c|} \hline 1 & -1 \\ \hline -1 & 1 \\ \hline \end{array} & \text{or } \begin{array}{|c|c|} \hline -1 & 1 \\ \hline 1 & -1 \\ \hline \end{array} & \text{if } T \text{ contains a} \\ & & \text{good switch point} \\ 0 & & \text{otherwise.} \end{cases}$$

Vertical Dividing Line: Passes through the geometric centre of  $T$

Slanting Dividing Line: Chosen such that the measure of balanced  $\alpha$  above and below the line are equal

Fact:  $\{G_i\}$  are orthogonal

$$\langle G_i, H \rangle = \sum \langle G_i, H_{a,b} \rangle$$



$$\langle G_i, H_{a,b} \rangle = 0$$



$$\langle G_i, H_{a,b} \rangle = 0$$

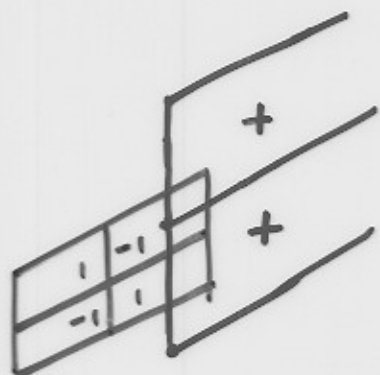


$$\langle G_i, H_{a,b} \rangle = 0$$

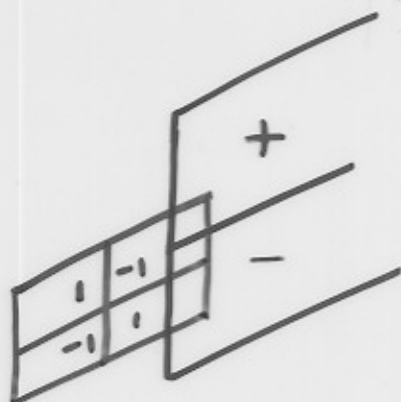


$$\langle G_i, H_{a,b} \rangle = 0$$

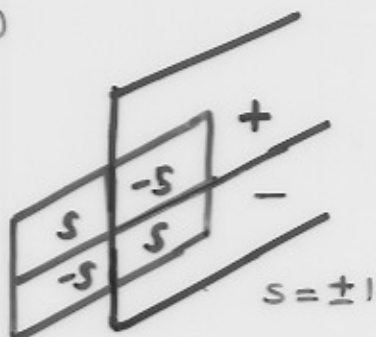




$$\langle G_L, H_{a,b} \rangle = 0$$



$$\langle G_L, H_{a,b} \rangle \neq 0$$



Switch point at the centre:  $\langle G_L, H_{a,b} \rangle \geq \frac{\mu_j^*}{8\beta}$

General position:  $E(\langle G_L, H_{a,b} \rangle) \geq \frac{(\mu_j^*)^2}{64\beta}$

$\mu_j^* = \frac{\mu_j}{\tau}$  where  $\mu_j$  is the measure of balanced  $\alpha$  in the  $j$ th sector

$$E(\langle G_L, H \rangle) \geq \frac{\sum (\mu_j^*)^2 s_j^*}{64\beta}$$

where  $s_j^*$  is the number of good switch points in the  $j$ th sector.

Question:  $\mu_j^*$  and  $s_j^*$  are large on average, but how well do they overlap?

Lemma: Let  $J \subseteq [0, 1)$  be an arbitrary interval of length  $\lambda$ , and let  $b_1, b_2, \dots, b_q$  be integers. Let  $N(\alpha, J) = |\{j : \{b_j \alpha\} \in J, 1 \leq j \leq q\}|$ . If  $q \geq \lambda^{-6}$ , then

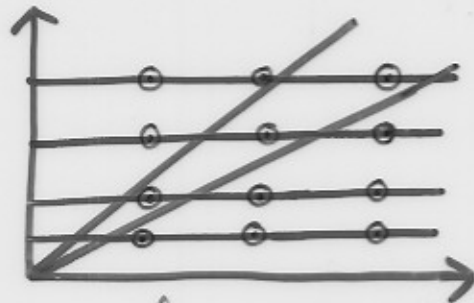
$$\mu(\alpha \in [0, 1] : N(\alpha, J) \geq 1) \geq 1 - \frac{1}{\sqrt{q}} \quad (\text{Beck '86})$$

Fact: Under the same hypotheses,

$$\mu(\alpha \in [0, 1] : N(\alpha, J) \geq \frac{q\lambda}{2}) \geq 1 - \frac{\delta}{\sqrt{q}}$$

Upshot:  $S_j^*$  is concentrated around its mean.

We confine our attention to "rich" sectors ( $\mu_j^* > \frac{\delta}{2}$ )



The measure of  $\alpha$  with fewer than half the expected number of switch points is small.

So "most" rich sectors contain "enough" switch points.

Can be shown: There are at most  $\frac{n^2}{4\beta}$  bad switch points.

(No lattice point: Typical; Two or more: Quite rare)

Putting it all together

$$E(\langle H, g_i \rangle) \geq \frac{c_1 n q}{\beta}$$

$$E(\langle H, g_i \rangle^2) \geq [E(\langle H, g_i \rangle)]^2 \geq \frac{c_1^2 n^2 q^2}{\beta^2}$$

$$\|g_i\| \leq \frac{n q}{2\beta}$$

$$E(\langle H, g_i \rangle^2) \geq \frac{c_1^2 n q}{2\beta} = \Omega\left(\frac{n^2}{(\log n)^{3/5}}\right)$$

$$\|H\|^2 \geq \frac{E\left(\sum_{i=1}^r \langle H, g_i \rangle^2\right)}{(3n^2/32)} = \Omega((\log n)^{3/5})$$

$\|H\| > (\log n)^{3/5}$  for sufficiently large  $n$ ; a contradiction.

Open Questions

- Better lower bound?
- Unbounded discrepancy for all but countably many  $\alpha$ ?
- Upper bound?